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Constructing and Analyzing the LSM Compaction Design Space





Log-Structured Merge-tree









tarantool







J. SCYLLΛ







































TM





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Why LSM?







fast writes



Why LSM ?

competitive reads



good space utilization



fast writes

competitive reads





good space utilization



fast writes

competitive reads







good space utilization



















workload





LSM tuning





COMPACTION

performance



Our Goal





Roadmap to pick compactions



Our Goal





Roadmap to pick compactions



Our Goal



Answer to complex design questions







Our Goal



learn from 2000+ experiments























































































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compaction



compaction?








































































How to organize the data on device?



How much data to move at-a-time?



Which block of data to be moved?



When to re-organize the data layout?





Compaction

granularity







Data movement policy













How to organize the data on device?

How much data to move at-a-time?

Which block of data to be moved?

When to re-organize the data layout?



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Data Layout number of runs per level







Data Layout number of runs per level



tiering [lazy]



1-leveling



leveling



Data Layout number of runs per level

L-leveling

tiering







Compaction Granularity





Compaction Granularity data moved per compaction









S Compaction Granularity data moved per compaction



2 Compaction Granularity data moved per compaction







2 Compaction Granularity data moved per compaction













S Compaction Granularity data moved per compaction













2 Compaction Granularity data moved per compaction











2 Compaction Granularity data moved per compaction











Data Movement Policy



files





Data Movement Policy which data to compact



files



Data Movement Policy which data to compact



coldest file





files



minimum overlap with parent level

file with most **tombstones**



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4 Compaction Trigger











level saturation







level saturation







level saturation

- number of sorted runs
- age of a file
- space amplification







Data Layout









Data Movement Policy

Compaction Trigger
Data Layout

Compaction Granularity



Any Compaction Algorithm



Data Movement Policy

Compaction Trigger

Database	Data layout	Compaction Trigger					Compaction Granularity				Data Movement Policy							
		Level saturation	#Sorted runs	File staleness	Space amp.	Tombstone-TTL	Level	Sorted run	File (single)	File (multiple)	Round-robin	Least overlap (+1)	Least overlap (+2)	Coldest file	Oldest file	Tombstone density	Expired TS-TTL	N/A (entire level)
RocksDB [30], Monkey [22]	Leveling / 1-Leveling	1		1					1	1		1		✓	1	1		
	Tiering		✓		✓	✓		✓										✓
LevelDB [32], Monkey (J.) [21]	Leveling	~							1		1	1	1					
SlimDB [47]	Tiering	✓							✓	✓								✓
Dostoevsky [23]	<i>L</i> -leveling	\checkmark^L	\checkmark^T				\checkmark^L	\checkmark^T				\checkmark^L						\checkmark^T
LSM-Bush [24]	Hybrid leveling	\checkmark^L	\checkmark^T				\checkmark^L	\checkmark^T				\checkmark^L						\checkmark^T
Lethe [51]	Leveling	1				✓			✓	✓		✓					✓	
Silk [11], Silk+ [12]	Leveling	1							✓	✓	✓							
HyperLevelDB [35]	Leveling	1							✓		✓	✓	✓					
PebblesDB [46]	Hybrid leveling	1							✓	✓								✓
Cassandra [8]	Tiering		✓	✓		✓		✓										✓
	Leveling	✓				✓			1	1		1				✓	✓	
WiredTiger [62]	Leveling	✓					✓											✓
X-Engine [34], Leaper [63]	Hybrid leveling	✓							✓	✓		✓				✓		
HBase [7]	Tiering		✓					1										✓
AsterixDB [3]	Leveling	✓					✓											✓
	Tiering		✓					✓										✓



10 compaction strategies



10 compaction strategies



10 compaction strategies





10 compaction strategies



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10 compaction strategies



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LSM tuning

Data moved (GB) 30 10







Data moved (GB)





Data moved (GB)





Data moved (GB)











Tiered data layout has the highest write throughput but also the highest tail write latency. Tail write latency (ms)

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Tiered data layout has the highest write throughput but also the highest tail write latency.





Compacting data at smaller granularity reduces data movement.

Tiered data layout has the highest write throughput but also the highest tail write latency.

Hybrid data layouts dominate point lookup performance.









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For update-intensive workloads, tiering dominates the performance space.





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Summary



Compaction as first-order design primitives.





Compaction as first-order design primitives.

Guidelines to design and tuning through experiments.









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Guidelines to design and tuning through experiments.







